

Differential Privacy is the dominant paradigm for designing learning algorithms that protect the privacy of <u>user data</u>.

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But we need a model of private learning that protects the privacy of the ML algorithm designer himself, too.

Learning with random examples

h: 20,13 ~ 20,13 $\Pr_{\mathbf{Y}}\left[\begin{array}{c} P_{\mathbf{F}}\\ \mathbf{Y} \\ \mathbf{Y$ (*1, E[*2)] (*1, E[*2)]



Agnostic Learning with random examples

h: 30,13 ~ 20,13 $\Pr_{\mathbf{Y}}\left[\begin{array}{c} P_{\mathbf{Y}}\\ P_{\mathbf{Y}}\\ \mathbf{Y} \\ \mathbf{Y} \\$ (*1, E(*2)) (*1, E(*2)) (*2, E(*2))



Agnostic Learning with random examples

h: 30,13 ~ 20,13 $\Pr_{\mathbf{Y}}\left[\begin{array}{c} P_{r}\\ P_{r}\\ \mathbf{Y} \sim \mathcal{U}\left[\begin{array}{c} h(\mathbf{x}) \neq g(\mathbf{x}) \\ \mathbf{Y} \end{array}\right] \leq \epsilon \geq 1-\delta$ (*1, E(*2)) (*1, E(*2)) (*2, E(*2)) Polynomial time.



Learning with queries

h: 20,13" -> 20,13 $\Pr_{f}\left[\begin{array}{c} P_{r} \\ x \sim \mathcal{U} \\ y \end{array}\right] \neq f(x) = f(x) = \frac{1}{2} = \frac{1}{2} = \frac{1}{2}$



Agnostic Learning with queries

h: 20,13" -> 20,13 $\begin{array}{c|c} P_r & P_r \\ \varphi & \times \sim \mathcal{U} \\ \uparrow & & & \\ \end{array} \begin{array}{c} h(x) \neq g(x) \\ \Rightarrow g(x) \\$



Agnostic Learning with queries

h: 20,13" ~> 20,13 $\begin{array}{c|c}
P_r & P_r \\
P_r & h(x) \neq g(x) \leq \varepsilon \geq 1-\delta \\
\chi \sim U & \chi \sim 0
\end{array}$ Polynomial time.



What's the difference (in polynomial time)?

$$\sum_{i \in S \subseteq [n]} \chi_i \pmod{2}$$



Non-noisy parity functions (Known: Gaussian elimination) Decision trees (unknown with random examples)

Agnostic Halfspace (Known: Kalai-Klivans-Mansour-Servedio, 2008)

 $\langle x, w \rangle$ + noise

Noisy parity functions (super unknown with random examples, see Learning Parity with Noise)

Queries are useful. So what do they reveal?

Kind of obvious, but let's get into it...

Learning with queries and an adversary

$$h: \frac{1}{2}0, 13^{n} \rightarrow \frac{1}{2}0, 13 \quad h \in \Lambda$$

$$\Pr\left[\Pr\left[\Pr\left(h(x) \neq f(x)\right] \le e\right] \ge 1 - \frac{1}{2}$$

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Simple noisy parity learning with queries





Simple noisy parity learning with queries and adversary





Adversary conducts the same majority vote to recover hidden vector.

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- So what do they reveal?

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- Queries are useful.
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- Sometimes, everything about the concept. Also, what the learner learned.
- Can we design algorithms with keyed query sets, that leak little information about the function, or the learner's intentions?
 - (To any <u>efficient</u> adversary, without knowledge of the key)

1. Prior knowledge used to influence the agnostic learning task remains private.

Consider the naive way of testing influence of variable of a function, or a choice of touchstone class in agnostic model

2. The concept itself remains unintelligible to anyone without the key to the query set. Desirable whenever the data labelling process is expensive or otherwise valuable

3. It can be desirable to obtain plausible deniability that any learning occurred at all.

. . .

Privacy Motivations?

$$h: \xi_0, i_3^{n} \rightarrow \xi_0, i_3^{n}$$

$$\Pr \left[\Pr_{x \sim \mathcal{U}} \left[h(x) \neq f(x) \right] \right]$$
Model stealing adversary extracts an approximation h to ML model f
Polynomial time

Model Stealing



Model Stealing with a defense

$$h: \frac{1}{2}0, 13^{n} \rightarrow \frac{1}{2}0, 13$$

$$Pr\left[Pr\left(h(x) \neq f(x)\right] \leq \epsilon\right] \geq 1-5$$

$$\frac{query}{house}$$

$$Polynomial time$$



Polynomial time

Model Stealing with a defense

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ML model owner **Polynomial time**

Example: Queries are "suspicious" when ML model owner can learn a good "proxy model" from them

E.g. the noisy parity setting we covered.

This has actually been proposed as a defense! See e.g. "Extraction Monitor" (Kesarwani-Mukhoty-Arya-Mehta, 2018)





Model Stealing with a defense from the adversary's perspective

$$h: \frac{1}{2}0, 13^{n} \rightarrow \frac{1}{2}0, 13$$

$$\Pr\left[\Pr\left(\frac{h(x) \neq f(x)}{x \sim u} \right) \leq \epsilon\right] \geq 1 - 5$$

$$\frac{quury}{hvelow}$$

$$\frac{quury}{hvelow}$$

Important: The model stealing adversary views the model owner as the "adversary's adversary."

Motivation: The model stealing adversary could bypass the model owner's defense if it could perform queries that "hide" what he is learning. See (<u>K</u>., 2023 for more)

See (K., 2023) for MUCH more info on Model stealing defenses



ML model owner **Polynomial time**

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Outsourcing of scientific experiments Drug Discovery; Quantitative Structure-Activity Research



"Experiments" are really just queries to a Boolean function (e.g. a specific set of molecules does or does not react).

There exists no protocol that a client can run with the nature. The results of the experiments are always viewed by a corrupt lab tech.

For various IP reasons, the client may want to privately run experiments (an "un-leakable" dataset, domain knowledge).

You may have noticed... We cannot hide everything (e.g

Main insight: hide only knowledge generated by queries.

Adversary learns no more than what is available by data occurring "in the wild." i.e., uniformly random examples.

We will not try to prevent leakage on any "too simple" class: that is, efficiently learnable with random examples.

We cannot hide everything (e.g. one of the two constant functions).





Defining Covert Learning (Canetti-K., 2021)

example distribution D, if there exists a polynomial time algorithm L satisfying:

and p.p.t. decision algorithm A,

$$\Pr_{A,L} \left[A \left(L_{\text{QuerySet}} \right) = 1 \right] -$$

In other words, the "transcript" of queries requested by L is computationally indistinguishable from random examples pulled from D

- A collection of hypothesis classes $C = \{H_i\}_{i \in [m]}$ is Covertly Learnable with respect to the
- 1. Agnostic Learning. For any $H_i \in C$, L with query access to f outputs h that satisfies:
 - $\Pr_{L^{f}(n,\epsilon,\delta,H)}\left[\Pr_{x\sim D}[h(x)\neq f(x)] \le \min_{h\in H}\Pr_{x\sim D}[\operatorname{opt}(H)(x)\neq f(x)] + \epsilon\right] \ge 1 \delta$
- 2. **Privacy.** There exists a p.p.t. simulation algorithm S such that, for any $H_i \in C, f$,
 - $-\Pr_{A,S}\left[A\left(S\left((X,f(X))\sim D\right)\right)=1\right]\right| \leq n^{-\omega(1)}$

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If agnostic learning over D is easy, then **Covert Learning is easy (by design).**



1. Hypothesis Hiding: prior knowledge used to influence the learning task remains private. Queries are indistinguishable whether you are learning polynomials dependent on one set of variables or another.

3. Undetectable: It can be desirable to obtain plausible deniability that any learning occurred at all. Nobody can decide that you applied a learning algorithm (Undetectable Model Stealing).

Benefits

2. Concept Hiding: the concept itself remains unintelligible to anyone without the key to the query set. Nobody can "free-ride" the learning process, if it is hard to learn with random examples.





Adversary conducts the same majority vote

Low noise LPN

Error bits are 1 with probability $n^{-1/2}$ and 0 otherwise

Secret bits are are 1 with probability $n^{-1/2}$ and 0 otherwise

Known to be as hard to find **w** as if it was uniformly random

Commonly assumed it takes $2^{n^{\epsilon}}$ time and random examples to find w.







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Make queries pseudorandom.

Apply one-time pads to these revealing queries

OTPs are pseudorandom by assumption that learning parities with noise rate $O(n^{-1/2})$ is hard

Observe this assumption is minimal

Simulator algorithm is immediate: output random examples given as input.





Sample random $n \times n$ matrix A

For every query $q_i \in \{0,1\}^n$ that we want to make,

choose a mask by $s_i A + v_i$ where $s_i, v_i \in \{0,1\}^n$ are sampled according to the noise distribution

The queries will be $q_i + s_i A + v_i$ and also the rows of A



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pseutorandom query
label

$$(s_iA + V_i + q_i) w + e_i + s_i Aw + \tilde{e})$$

 $V_iW + q_iW + e_i$
 $=$
 $\tilde{E} V_{ij}W_j = 1] = \frac{1}{h}$
 $\therefore Pr[V_{ij}W_j = 1] \le 0.449$
for sufficiently large n
 $Pr[S_i\tilde{e} = 1] \le 0.49$
for sufficiently large n
Covert Learning parities with low noise over Unif (Canetti-K., 2021)

pseudorandom query
label

$$(s_iA + V_i + q_i) w + e$$

 $V_iw + q_iw$
 $=$
 $V_{ij}w_j = 1] = \frac{1}{n}$
 $\therefore Pr[V_{ij}w_j = 1] \le 0.49$
for sufficiently large r





Covert Learning parities with low noise over Unif (Canetti-K., 2021)

Summary

1. Given secret key, we are able to make queries on a noisy parity function $w \in \{0,1\}^n$ Enough to learn!

2. If n is very large, you may use prior knowledge that all relevant points exist in a subset of size k. Naive method reveals subset. You may now hide subset with n + O(k) queries. Naive hiding method uses O(n) queries.

3. $2^{n^{\epsilon}}$ -time adversary that obtains query set learns little about w, unless LPN is false.

"Local" model (Jawale-Holmgren, 2023)



Clearly very relevant still to model stealing, outsourcing lab experiments.

Consider a "sybil" attack for model stealing, or 2 noncolluding science labs.





A simple algorithm for k-juntas in the "Local" model Implicit in (Ishai-Kushilevitz-Ostrovsky-Sahai, 2019) and (Canetti-K., 2021)

On a function $f: \{0,1\}^n \rightarrow \{0,1\}$, a variable x_i is irrelevant if for any input x, f(x) = f(y) for y = xexcept at the i^{th} bit. A variable is <u>relevant</u> iff it is not irrelevant.

 $f: \{0,1\}^n \rightarrow \{0,1\}$ is a k-junta if it has at least n - k irrelevant variables.

Best known algorithms for learning k-juntas with uniformly random examples go in $n^{\epsilon k}$ time for some $\epsilon > 2/3$.

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(viewed independent of right side)

the Covert Learning simulation privacy definition.



A simple algorithm for k-juntas in the "Local" model

Implicit in (Ishai-Kushilevitz-Ostrovsky-Sahai, 2019) and (Canetti-K., 2021)

With probability $\approx 2^{-k}$, we obtained $\langle x_i, f(x_i) \rangle$ which is sensitive at x_j , and we also queried for a neighbor which at a relevant index.

> Uniformly random neighbors \leftarrow of each x_i

With $2^{O(k)}$ neighbor/random example pairs, we can identify all relevant variables.

With $2^{O(k)}$ more random examples, we can find	Wh
and memorize the exact truth table.	but



hen $k = O(\log n)$, this is polynomial time for the learner, it quasipolynomial time for the adversary, by simulation.

Future directions

We know globally Covert Learning for low degree Fourier coefficients and decision trees.

Can every query algorithm be compiled into a (global) covert learning algorithm, under the (minimal) assumption that the task is hard with random examples?

- Standard LPN?

- AC0[p]? (Carmosino-Impagliazzo-Kabanets-Kolokolova, 2016)

- Covert Learning for non-Boolean functions? Maybe use CLWE? (Bruna-Regev-Song-Tang, 2021)

Zou et al., 2023

Should AI model respond to client?

Worth exploring: AI Jailbreaking

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Al needs to predict, given queries (and what it knows about the world), whether client will compute something it isn't supposed to.



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Al needs to predict, given queries (and what it knows about the world), whether client will compute something it isn't supposed to.

If Queries are distributed identically to the collected data, then Al fundamentally cannot decide whether it should respond or not.



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"The only winning move is not to play"







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